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ON THE LOCAL AND GLOBAL OPTIMUMS IN DISTRIBUTION FEEDER RECONFIGURATION

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A simple algorithm, based on iterative stepwise shifting of loads (switch exchange) between distribution network feeders, is presented. The results of a computer experiment in local distribution network are presented and compared with the results of other researches.

Introduction

The scope of one feeder in the radial distribution network is determined by breakpoints (lines with switches open) separating this feeder from the rest of the network. Depending on cable lengths, their cross-sections and customers' loads, different feeder configurations give different total losses in the distribution network. Reduction of the losses might be significant, and therefore the problem of optimal feeder configuration has been one of the actual problems in the research of distribution networks for a long time.

The present development of SCADA systems and remote control facilities make the problem actual. Various methods are applied for optimizing feeder configuration: simple exchange of tie and sectionalizing of switches [5], more sophisticated sequential switch opening [10], network partitioning [11], application of generic algorithms [17], reduced-order binary decision diagrams [19]. Simple switch exchange methods are considered inefficient as they often lead to local optimality "traps", and for large networks, only methods based on heuristics and on combinatory methods are considered suitable.

Nevertheless, the simplest approach for reconfiguration of distribution feeder is used in the algorithm presented below. It is found to be usable, fast, and it has given a considerable reduction of losses. It was found that different switching order produces different optimal network configurations. Therefore, the question about local and global optimality remains open.

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Fig. 1. A breakpoint between feeders F1 and F2

An Iterative Method of Shifting of Breakpoints

The method presented here is a modification of classical switch exchange method. The objective of the algorithm is to determine the optimal location of each breakpoint individually. Therefore, as each breakpoint is considered separately, the problems like the 2^n or similar problems of combinatory analysis do not occur. Nevertheless, network losses can be not the only criteria considered in such approach.

Suppose the initial configuration of a distribution network with radial structure is given, where no normal constraints are violated, all line currents and bus voltages are within normal limits. As a simplification, the loads are modeled as constant currents. The following algorithm will be applied to this network:

1. Load-flow calculation is made for the whole network in order to determine the voltages of the bus-bars of supply substations

$$U_{s} = U_{s}^{a} + jU_{s}^{r}, s = 1, 2, ... S$$

where S – the set of supply substations.

- 2. A list of all breakpoints will be composed. The way of its composition and the order of the breakpoints in the list are free. Suppose the number of breakpoints is *B*, breakpoints numbers are N_b , b=1, 2, 3, ..., B, respectively. Additional information on every breakpoint and every switch is needed whether it is permitted to close this breakpoint (switch) a. The same applies to every line switch whether it is permitted to switch this line off (usual operational limitations should be taken into consideration, if those exist).
- 3. Beginning with the first one in this list the following is calculated:
 - 3.1. Breakpoint will be checked for the possibility of its shifting. If the closure of this switch is not permitted, the next breakpoint in the list will be studied.
 - 3.2. Two vectors of lines will be composed: one from one side of the breakpoint to the supply bus *s*' (feeder *F1*), and second from the other side of the breakpoint to the supply bus *s*'' (feeder *F2*) (i.e. circuits from the breakpoint to the supply buses (Fig.1).

Suppose the first vector (circuit) having *N* elements (n=1, 2, ...N) and the second one *M* elements (m=1, 2, ...M). The elements of those vectors (lines) are numbered starting from supply buses.

3.3. Calculate voltages or voltage drops on both sides of the breakpoint

$$U_{1} = U_{s'} - \sum_{n=1}^{N} (I_{n}^{*} Z_{n}^{*}), \forall (n, n = 1, ...N),$$

$$U_{2} = U_{s''} - \sum_{m=1}^{M} (I_{m}^{*} Z_{m}^{*}), \forall (m, m = 1...M)$$
(1)

3.4. Compare voltages U_1 and U_2 . Relation $\begin{vmatrix} \bullet \\ U_1 \end{vmatrix} > \begin{vmatrix} \bullet \\ U_2 \end{vmatrix}$ indicates that the

voltage drop in feeder F1 is less than in feeder F2. It could indicate that feeder F2 is carrying more load than feeder F1, and there could be some reason to switch some load from feeder F2 to feeder F1, closing the breakpoints switch and opening the switch of line number M, thus transferring load I_M to feeder F1.

- 3.5. Check if the switching off of line number M is permitted. If it is, then proceed with the algorithm, if not go to the next breakpoint.
- 3.6. Calculate the losses in the same circuits where the voltage drop was calculated

$$\Delta P_{1}' = \sum_{n=1}^{N} \{ (I_{a,n}^{2} + I_{r,n}^{2}) * r_{n} \}, \Delta P_{2}' = \sum_{m=1}^{M} \{ (I_{a,m}^{2} + I_{r,m}^{2}) * r_{m} \} \Delta P_{\Sigma}' = \Delta P_{1}' + \Delta P_{2}' .$$
⁽²⁾

3.7. Calculate losses in the same circuits with load I_M switched from feeder F1 to feeder F2

$$\Delta P_{1}^{"} = \sum_{n=1}^{N} \left[(I_{a,n} + I_{a,M})^{2} + (I_{r,n} + I_{r,M})^{2} \right] * r_{n} + (I_{a,M}^{2} + I_{r,M}^{2}) * r_{0}$$

$$\Delta P_{2}^{"} = \sum_{m=1}^{N} \left[(I_{a,m} - I_{a,M})^{2} + (I_{r,m} - I_{r,M})^{2} \right] * r_{m}$$

$$\Delta P_{\Sigma}^{"} = \Delta P_{1}^{"} + \Delta P_{2}^{"}$$
(3)

3.8. Calculate possible change of losses by shifting of the breakpoint

$$\Delta \Delta P = \Delta P_{\Sigma}' - \Delta P_{\Sigma}'' \tag{4}$$

3.9. If $\Delta\Delta P > 0$, the shifting of the breakpoint or the load I_M from feeder *F1* to feeder *F2* results in loss reduction.

If the shapes of the load curves of different buses are sharply different, the energy losses might be considered instead of power losses.

The levels of line currents and bus voltages are checked. If they are within admissible limits, the new breakpoint will be fixed, topology of the network will be changed, the old breakpoint in the list of breakpoints will be replaced with the new one.

3.10. Go to the next breakpoint in the list.

4. After first inspection of all breakpoints, the next iteration begins. All breakpoints are tested in the same order. The iterations are repeated until no breakpoint could be shifted during whole iteration.

5. Change the order of testing the breakpoints in order to find a possible more optimal configuration, and start the whole procedure from the beginning.

Computer Experiments in Local Distribution Network

For testing the applicability and efficiency of the developed algorithm, calculation was made with actual data of a part of the local distribution network. The results of this computer experiment are given in Table 1.

Different paths shown here have the following meaning:

Path 1 – the breakpoints are handled in the order of their listing.

Path 2 – the testing of breakpoints begins from the end of their list.

Path 3 – breakpoints are taken alternatively from the beginning and from the end of the breakpoints list.

Path 6 – the breakpoints are sorted by the possible loss reduction that could be obtained by their shifting.

Then, beginning from the breakpoint with the maximal loss reduction down to the end of the list the shifting process is executed. The difference of the loss reductions for different order of breakpoints is 0.6% for the first network and 2.1% for the second network. Bearing in mind that all bus loads are given with an accuracy of about 5%, these differences could be considered acceptable. Reduction of losses by around 10% could also be considered essential.

	North-east districts (10 kV)		South-west districts (6k			(6kV)		
Number of supply substations	3		7					
Number of feeders	42		83					
Number of load buses	522		555					
Number of breakpoints	237		172					
Number of fixed breakpoints	22		83					
Total load in the network, MW	63.33		79.64					
Total line losses, initial network, kW	518.0		579.8					
Number of the path (test order)	1 2 3 6		1	2	3	6		
Iterations calculated	5	5	5	10	4	4	4	10
Shifted breakpoints	41	39	36	42	41	42	44	44
Line losses after optimisation, kW	434.7	435.8	438.0	436.4	518.7	508.7	508.3	520.7
Reduction of losses, kW	83.4	82.1	80.1	81.6	61.1	71.0	71.4	59.1
Reduction of losses, %	16.1	15.9	15.5	15.8	10.5	12.3	12.3	10.2

Table 1. Search for Optimal Feeder Configuration in Local City Network

Nevertheless, the question remains – how far is the best solution found here from the global optimum. To answer this question a new algorithm using GA or similar heuristics had to be developed. Another idea for testing the algorithm developed here is to use this algorithm against the work of other authors; using the data of their test networks to compare our results with the results of algorithms developed by other authors. In all papers listed in the reference [1-16], the test networks were too small, or the data given in the paper was incomplete. A comparison possibility appeared in 2002, when J. Z. Zhu published his paper in the journal "Electric Power System Research" [17], giving complete data of the test network he had used.

A Comparison with the Results of Zhu [17]

In the fifth chapter of his paper, J. Z Zhu tests his approach to the feeder configuration optimization on 16-bus and 33-bus distribution systems. We consider here the 33-bus network. The configuration of this system is given in Fig. 2, the data and parameters are brought in Table 2.

Table 2. Parameters and Loads of the Network in Fig. 2

	1					r		r
Line No.	Node I	Node J	R [Ω]	$Xl\left[\Omega\right]$	$P_{j}[kW]$	Q_j [kVar]	$V_j Mod p.u.$	$V_j \Delta$ [rad]
1	1	2	0.0922	0.047	100.00	60.00	0.997	0.0145
2	2	3	0.493	0.2512	90.00	40.00	0.9829	0.096
3	3	4	0.3661	0.1864	120.00	80.00	0.975	0.0617
4	4	5	0.3811	0.1941	60.00	30.00	0.9681	0.2283
5	5	6	0.819	0.707	60.00	20.00	0.9497	0.1339
6	6	7	0.1872	0.6188	200.00	100.00	0.9462	-0.0964
7	7	8	0.7115	0.2351	200.00	100.00	0.9413	-0.0603
8	8	9	1.0299	0.74	60.00	20.00	0.9351	-0.1334
9	9	10	1.044	0.74	60.00	20.00	0.9292	-0.1959
10	10	11	0.1967	0.0651	45.00	30.00	0.9284	-0.1887
11	11	12	0.3744	0.1298	60.00	35.00	0.9269	-0.1785
12	12	13	1.468	1.1549	60.00	35.00	0.9208	-0.2698
13	13	14	0.5416	0.7129	120.00	80.00	0.9185	-0.3485
14	14	15	0.5909	0.526	60.00	10.00	0.9171	-0.3862
15	15	16	0.7462	0.5449	60.00	20.00	0.9157	-0.4094
16	16	17	1.2889	1.721	60.00	20.00	0.9137	-0.4868
17	17	18	0.732	0.5739	90.00	40.00	0.9131	-0.4963
18	2	19	0.164	0.1565	90.00	40.00	0.9965	-0.0037
19	19	20	1.5042	1.3555	90.00	40.00	0.9929	-0.0633
20	20	21	0.4095	0.4684	90.00	40.00	0.9922	-0.0827
21	21	22	0.7089	0.9373	90.00	40.00	0.9916	-0.103
22	3	23	0.4512	0.3084	90.00	50.00	0.9794	-0.065
23	23	24	0.898	0.7091	420.00	200.00	0.9727	-0.0237
24	24	25	0.8959	9.7071	420.00	200.00	0.9694	-0.0674
25	6	26	0.2031	0.1034	60.00	25.00	0.9477	0.1734
26	26	27	0.2842	0.1447	60.00	25.00	0.9452	0.2295
27	27	28	1.0589	0.9338	60.00	20.00	0.9337	0.3124
28	28	29	0.8043	0.7006	120.00	70.00	0.9255	0.3904
29	29	30	0.5074	0.2585	200.00	100.00	0.9219	0.4956
30	30	31	0.9745	0.9629	150.00	70.00	0.9178	0.4112
31	31	32	0.3105	0.3619	210.00	100.00	0.9169	0.3882
32	32	33	0.3411	0.5302	60.00	40.00	0.9166	0.3805
33	25	29	0.5	0.5				
34	8	21	2	2				
35	12	22	2	2				
36	9	15	2	2				
37	18	33	0.5	0.5				



Fig. 2. 33-bus test system in [17]

Table 3. Optimization of the 33-Bus Test Network by the Method Presented in This Paper

Break-point No.Existing open switch No.Possible loss reduction, WNew open switch No.Real loss reduction, W13316290281629023414850714850335132111113211436114014110253714032199Iteration No.2128241033243527-192707031191010885414-3170140532-620320Iteration No.3133-2464033027-17090703106809699414-5040140532-820320Iteration No.4133-246403327-17090703106809699414-5040140532-820320Iteration No.4133-246403327-176807039-637090414-5040140532-820320Total loss reduction 49,670 W <th></th> <th>-</th> <th></th> <th></th> <th></th>		-						
Iteration No.1Iteration No.113316290281629023414850714850335132111113211436114014110253714032199Iteration No.212824103327-192707031191010885414-3170140532-620320Iteration No.3133-2464033027-17090703106809699414-5040140532-820320Iteration No.4133-246403327-17090703106809699414-5040140532-820320Iteration No.4133-2464033027-176807039-637090414-5040140532-820320	Break-point		Possible					
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	No.	open switch No.	loss reduction, W	open switch No.	loss reduction, W			
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	Iteration No.1							
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	1	33	16290	28	16290			
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	2	34	14850					
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	3	35	13211	11	13211			
Iteration No.2128241033243527-192707031191010885414-3170140532-620320Iteration No.3133-2464033027-17090703106809699414-5040140532-820320Iteration No.4133-24640739-637090414-5040140532-820320		36	1140	14	1102			
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	5	37	140	32	199			
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	Iteration No.2							
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	1	28	2410	33	2435			
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	2	7	-19270	7	0			
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	3	11	910	10	885			
Iteration No.3133 -24640 33027 -17090 703106809699414 -5040 140532 -820 320Iteration No.4133 -24640 3327 -17680 7039 -6370 90414 -5040 140532 -820 320	4	14	-3170	14	0			
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	5	32	-620	32	0			
$\begin{array}{c c c c c c c c c c c c c c c c c c c $	Iteration No.3							
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	1	33	-24640	33	0			
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	2	7	-17090	7	0			
5 32 -820 32 0 Iteration No.4 1 33 -24640 33 0 2 7 -17680 7 0 3 9 -6370 9 0 4 14 -5040 14 0 5 32 -820 32 0	3	10	680	9	699			
Iteration No.4133 -24640 33027 -17680 7039 -6370 90414 -5040 140532 -820 320	4	14	-5040	14	0			
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	5	32	-820	32	0			
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	Iteration No.4							
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	1	33	-24640	33	0			
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	2	7	-17680	7	0			
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	3	9	-6370	9	0			
lll	4	14	-5040	14	0			
Total loss reduction 49,670 W	5	32	-820	32	0			
	Total loss reduction 49,670 W							

The 33-bus test system consists of one supply transformer and 32 load points (distribution transformers). There are five initially open switches in the test system – "33", "34", "35", "36" and "37". The total system load is 3.715 MW, the basis of the system is V=12.66 kV and S=10 MVA.

Applying his method to the network presented in Fig. 2 and Table 2, J. Z. Zhu finds optimal configuration of this network by given loads with open switches No. 7, 9, 14, 32 and 33. Applying the method proposed in [6] to the same network, J. Z. Zhu finds optimal configuration to be with open switches No. 7, 10, 14, 33 and 37 and total losses by this configuration being about 10%

more than by the configuration found by his method. So, J. Z. Zhu concludes that global optimum is found in his paper and not in [6].

Two calculations were made for the network presented in Fig. 2 and Table 2. In order to get some comparison with the efficiency of the algorithm developed in this paper, in the first case the initial state of the network was taken with open switches No. 33, 34, 35, 36 and 37 – the initial state of the network. The path and results of this calculation are presented in Table 2.

As one can see, after the fourth iteration the configuration of networks has converged to the configuration obtained by the method published in [17].

In the second case the optimization result by the method published in [6] was taken for the initial state – switches No. 7, 10, 14, 33 and 37 are open. The results are given in Table 4.

The applied algorithm has found by the first step, two possible ways for loss reduction in the network configuration, which was considered as optimal in [6].

Break-point No.	Existing open switch No.	Possible loss reduction, W	New open switch No.	Real loss reduction, W				
	Iteration No. 1							
1	7	-17110	7	0				
2	10	480	9	503				
3	14	-4010	14	0				
4	33	-13540	33	0				
5	37	800	32	851				
	Iteration No. 2							
1	7	-14680	7	0				
2	9	-6370	9	0				
3	14	-5040	14	0				
4	33	-11710	33	0				
5	32	-820	32	0				
	Total loss reduction 1,354 W							

Table 4. Optimization of the "Optimized" by the Method of [6] Network

Conclusions

A variation of the classical switch exchange method is presented in this paper. The developed algorithm and software are fast and enable online control of large networks with considerable reduction of network losses. A comparison with the method using GA [17] is given showing on the test network similarity. However, it will be inconsiderate to declare that the algorithm presented here gives always the same results as methods based on GA. It is obvious that simple networks are not suitable for comparing different optimization methods. Decision on the suitability of control and optimization methods should be made based on the experiments in actual networks. More comparisons of different methods for optimization of configuration of distribution networks are needed. In addition, one more question needs to be

answered – how does the accuracy of the load data used by optimization influence the accuracy of optimum determination. Does the search for global optimum have a meaning at all, or are the local optimums found here sufficient for practical applications?

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